

# Modeling and Fabrication with Specified Discrete Equivalence Classes - Supplement

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## ACM Reference Format:

Zhong-Yuan Liu, Zhan Zhang, Di Zhang, Chunyang Ye, Ligang Liu, and Xiao-Ming Fu. 2021. Modeling and Fabrication with Specified Discrete Equivalence Classes - Supplement. 1, 1 (April 2021), 3 pages. <https://doi.org/10.1145/nnnnnnn.nnnnnnn>

## 1 APPENDIX

### 1.1 Proposition 1

*Prop. 1.* The shape of  $\Delta u_0 u_1 u_2$  is independent of the translation  $\mathbf{b}$ , and so do  $C$ . Given any  $R$ , the best  $\mathbf{b}$  minimizing  $\|\xi(R)\|_\infty$  is the vector from the origin  $\mathbf{o}$  to the center  $\mathbf{c}$  of  $C$ , and the minimum  $\|\xi(R)\|_\infty$  is the radius  $r_c$  of  $C$ .

**PROOF.** Independence of the translation is deduced since the edge vector  $\mathbf{u}_j \mathbf{u}_i = \mathbf{u}_i - \mathbf{u}_j = (\mathbf{v}_i - R\mathbf{p}_i - \mathbf{b}) - (\mathbf{v}_j - R\mathbf{p}_j - \mathbf{b}) = (\mathbf{v}_i - R\mathbf{p}_i) - (\mathbf{v}_j - R\mathbf{p}_j)$  is independent of the translation  $\mathbf{b}$ . Hence, the shape of  $\Delta u_0 u_1 u_2$  and the minimum covering circle (MCC)  $C$  of  $\Delta u_0 u_1 u_2$  are independent of the translation  $\mathbf{b}$  (Fig. 1).

Given a rotation  $R$ , we get the rotated template triangle  $\Delta p'_0 p'_1 p'_2$ , where  $\mathbf{p}'_i = R\mathbf{p}_i$ .  $\Delta u'_0 u'_1 u'_2$  is the triangle without the translation  $\mathbf{b}$ , i.e.,  $\mathbf{u}'_i = \mathbf{p}'_i - \mathbf{v}_i$ . We denote  $V_{\Delta u'_0 u'_1 u'_2} = \{\mathbf{u}'_0, \mathbf{u}'_1, \mathbf{u}'_2\}$ . Let  $\mathbf{q} = \mathbf{b} + \mathbf{o}$ . Define  $d_S(\mathbf{v}) = \max_{s \in S} \|\mathbf{v} - \mathbf{s}\|_2$  as the distance from a point  $\mathbf{v}$  to a point set  $S$ . Then, we have:

$$\begin{aligned} \|\xi\|_\infty &= \max\{\|\mathbf{v}_0 - \mathbf{p}'_0 - \mathbf{b}\|_2, \|\mathbf{v}_1 - \mathbf{p}'_1 - \mathbf{b}\|_2, \|\mathbf{v}_2 - \mathbf{p}'_2 - \mathbf{b}\|_2\} \\ &= \max\{\|\mathbf{u}'_0 - \mathbf{b}\|_2, \|\mathbf{u}'_1 - \mathbf{b}\|_2, \|\mathbf{u}'_2 - \mathbf{b}\|_2\} \\ &= \max\{\|\mathbf{u}'_0 - \mathbf{q}\|_2, \|\mathbf{u}'_1 - \mathbf{q}\|_2, \|\mathbf{u}'_2 - \mathbf{q}\|_2\} \\ &= d_{V_{\Delta u'_0 u'_1 u'_2}}(\mathbf{q}). \end{aligned}$$

The problem  $\min \|\xi\|_\infty$  is converted to finding the best point  $\mathbf{q}$  to minimize the distance from  $\mathbf{q}$  to the vertices of  $\Delta u'_0 u'_1 u'_2$ .

We claim that the best point  $\mathbf{q}$  for minimizing  $d_{V_{\Delta u'_0 u'_1 u'_2}}(\mathbf{q})$  is the center  $\mathbf{c}$  of  $C$ . Otherwise, there must be a point  $\mathbf{o}^*$  s.t.  $d_{V_{\Delta u'_0 u'_1 u'_2}}(\mathbf{o}^*) < d_{V_{\Delta u'_0 u'_1 u'_2}}(\mathbf{c}) = r_c$ . Then, setting  $r^* = d_{V_{\Delta u'_0 u'_1 u'_2}}(\mathbf{o}^*)$ , the circle  $C_{\mathbf{o}^*}(r^*)$  covers all the vertices of  $\Delta u'_0 u'_1 u'_2$  and  $r^* < r_c$ , which contradicts

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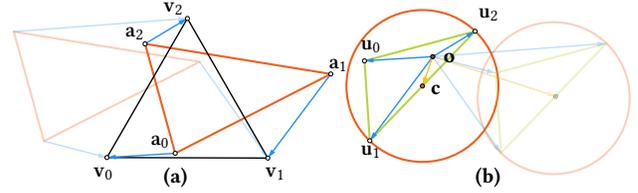


Fig. 1. Independent of the translation. (a)  $\Delta v_0 v_1 v_2$  is a triangle of the remeshed mesh  $\mathcal{R}$  and  $\Delta a_0 a_1 a_2$  indicates the template triangle  $\Delta p_0 p_1 p_2$  after a rigid transformation  $R$ . (b) Move the starting points of the three vectors  $(\mathbf{v}_0 - \mathbf{a}_0, \mathbf{v}_1 - \mathbf{a}_1, \mathbf{v}_2 - \mathbf{a}_2)$  to the origin  $\mathbf{o}$ . The orange circle is the MCC of  $\Delta u_0 u_1 u_2$  and  $\mathbf{c}$  is its center. Here  $\Delta u_0 u_1 u_2$  is an obtuse triangle and  $\mathbf{c}$  is the midpoint of its longest edge. The transparent figures show another case with the same rotation but the different translation. The shape of green triangle  $\Delta u_0 u_1 u_2$  remains the same, so do the MCC.

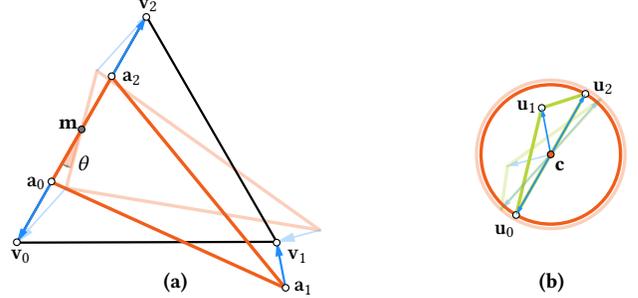


Fig. 2. The obtuse triangle case. (a)  $\Delta v_0 v_1 v_2$  is a triangle of  $\mathcal{R}$  and  $\Delta a_0 a_1 a_2$  indicates the template triangle  $\Delta p_0 p_1 p_2$  after best rigid transformation  $R$ .  $\mathbf{b}$ .  $\mathbf{a}_0 \mathbf{a}_2$  and  $\mathbf{v}_0 \mathbf{v}_2$  coincide and their midpoints coincide at a point, denoted as  $\mathbf{m}$ . (b) Move the starting points of the three vectors  $(\mathbf{v}_0 - \mathbf{a}_0, \mathbf{v}_1 - \mathbf{a}_1, \mathbf{v}_2 - \mathbf{a}_2)$  to the origin  $\mathbf{o}$ . The orange circle is the MCC of  $\Delta u_0 u_1 u_2$  and  $\mathbf{c}$  is its center.  $\mathbf{u}_0 \mathbf{u}_2$  is the longest edge. Based on *Prop. 1*, the origin  $\mathbf{o}$  is at the center  $\mathbf{c}$  of  $C$  after the best translation  $\mathbf{b}$ . The transparent figures show another case with the best translation but another rotation. Corresponding edges  $\mathbf{e}_f$  and  $\mathbf{e}_t$  do not coincide and the radius of its MCC is bigger.

the definition of MCC. As a result,  $\mathbf{q} = \mathbf{c}$  and  $\mathbf{b} = \mathbf{c} - \mathbf{o}$ . Since the MCC of  $\Delta u_0 u_1 u_2$  is independent of the translation  $\mathbf{b}$ ,  $\min \|\xi\|_\infty = d_{V_{\Delta u_0 u_1 u_2}}(\mathbf{c}) = r_c$ .  $\square$

### 1.2 Proposition 2

*Prop. 2.* The longest edge of  $\Delta u_0 u_1 u_2$  corresponds to an edge of  $\mathbf{f}$  (denoted as  $\mathbf{e}_f$ ) and an edge of  $\mathbf{t}$  (denoted as  $\mathbf{e}_t$ ), respectively. Then, if  $\Delta u_0 u_1 u_2$  is an obtuse triangle when  $\|\xi\|_\infty$  reaches the minimum, then  $\mathbf{e}_f$  and  $\mathbf{e}_t$  coincide and their midpoints coincide.

**PROOF.** Without the loss of generality, assume the longest edge is  $\mathbf{u}_0 \mathbf{u}_2$  and the best rotation is  $R^*$  when  $\|\xi\|_\infty$  obtains the minimum. Considering MCC in the obtuse case, the radius  $r_c$  is the half of the longest side  $\mathbf{u}_0 \mathbf{u}_2$  and the center  $\mathbf{c}$  is at the midpoint of  $\mathbf{u}_0 \mathbf{u}_2$ . Based on *Prop. 1*, the origin  $\mathbf{o}$  is at the center  $\mathbf{c}$  of  $C$  after the best

translation  $\mathbf{b}$ , indicating that  $\overrightarrow{a_0 v_0} = \overrightarrow{c u_0} = -\overrightarrow{c u_1} = -\overrightarrow{a_2 v_2}$ . Hence, the midpoints of  $a_0 a_2$  and  $u_0 u_2$  coincide at a point, denoted as  $\mathbf{m}$  (Fig. 2).

Next, we prove that  $\mathbf{e}_f$  and  $\mathbf{e}_t$  coincide by contradiction. Since  $\Delta u_0 u_1 u_2$  and  $C$  are independent of the translation, the center of rotation does not influence the shape of  $\Delta u_0 u_1 u_2$  and  $C$ . Without the loss of generality, we now rotate  $\Delta a_0 a_1 a_2$  around  $\mathbf{m}$ . Suppose that  $a_0 a_2$  and  $u_0 u_2$  are not coincided after the best rotation  $R^*$ . Since  $\Delta u_0 u_1 u_2$  is an obtuse triangle, we have  $|a_1 v_1| = |c u_1| < |c u_0| = |a_0 v_0| = |a_2 v_2| = r$ . Thus, we can rotate  $\Delta a_0 a_1 a_2$  a small angle  $\delta_\theta$  to a new rotated template triangle, denoted as  $\Delta a'_0 a'_1 a'_2$ , whose  $|v_1 a'_1|$  is still smaller than  $|v_0 a'_0| = |v_2 a'_2| = r'$  and  $r' < r$ . However,  $r' < r$  contradicts the assertion that rotation  $R^*$  is the best rotation.  $\square$

### 1.3 $f(\theta)$

Based on Prop. 1, the shape of  $\Delta u_0 u_1 u_2$  and  $C$  are also independent of the center of rotation. In the acute triangle case, we place  $a_0$  at the origin  $\mathbf{o}$ , then the center of rotation  $R$  is  $a_0$ . We draw the auxiliary lines, as shown in Fig. 3, where quadrilaterals  $a_0 a_1 v_1 u_1$  and  $a_0 a_2 v_2 u_2$  are parallelograms. Then,  $\Delta v_0 u_1 u_2$  is the same as  $\Delta u_0 u_1 u_2$ . Thus, we only need to find the radius of the circumcircle of  $\Delta v_0 u_1 u_2$ :

$$\min_R r_c = \min_R \frac{\|u_1 - v_0\|_2 \cdot \|u_2 - u_1\|_2 \cdot \|v_0 - u_2\|_2}{4 \text{Area}(\Delta v_0 u_1 u_2)},$$

where

$$R = \begin{bmatrix} \cos(\theta) & -\sin(\theta) \\ \sin(\theta) & \cos(\theta) \end{bmatrix}.$$

The expression for  $r_c$  can be derived as:

$$\begin{aligned} f(\theta) = r_c &= \frac{\|u_1 - v_0\|_2 \cdot \|u_2 - u_1\|_2 \cdot \|v_0 - u_2\|_2}{4 \text{Area}(\Delta v_0 u_1 u_2)} \\ &= \sqrt{\frac{(\alpha_1 + \alpha_2 \cos \theta + \alpha_3 \sin \theta)(\alpha_4 + \alpha_5 \cos \theta + \alpha_6 \sin \theta)(\alpha_7 + \alpha_8 \cos \theta + \alpha_9 \sin \theta)}{(\alpha_{10} + \alpha_{11} \cos \theta + \alpha_{12} \sin \theta)^2}} \\ &= \sqrt{\frac{(\alpha_1 + \alpha_2 \frac{1-t^2}{t^2+1} + \alpha_3 \frac{2t}{t^2+1})(\alpha_4 + \alpha_5 \frac{1-t^2}{t^2+1} + \alpha_6 \frac{2t}{t^2+1})(\alpha_7 + \alpha_8 \frac{1-t^2}{t^2+1} + \alpha_9 \frac{2t}{t^2+1})}{(\alpha_{10} + \alpha_{11} \frac{1-t^2}{t^2+1} + \alpha_{12} \frac{2t}{t^2+1})^2}} \\ &= g(t), \end{aligned}$$

where  $t = \tan(\theta/2)$  and  $\alpha_1, \dots, \alpha_{12}$  are the constants related to triangles  $\mathbf{f} = \Delta v_0 v_1 v_2$  and  $\mathbf{t} = \Delta p_0 p_1 p_2$ ,

$$\begin{aligned} \alpha_1 &= x_0^2 - 2x_0 x_{v_1} + x_{v_1}^2 + x_0^2 - 2x_0 x_{p_1} + x_{p_1}^2 + y_0^2 - 2y_0 y_{v_1} + y_{v_1}^2 + y_0^2 - 2y_0 y_{p_1} + y_{p_1}^2, \\ \alpha_2 &= 2(x_{v_1}(x_0 - x_{p_1}) + x_{v_0}(-x_0 + x_{p_1}) - (y_0 - y_{v_1})(y_0 - y_{p_1})), \\ \alpha_3 &= 2(x_{p_1}(y_0 - y_{v_1}) + x_{p_0}(-y_0 + y_{v_1}) + (x_{v_0} - x_{v_1})(y_0 - y_{p_1})), \\ \alpha_4 &= x_0^2 - 2x_0 x_{v_2} + x_{v_2}^2 + x_0^2 - 2x_0 x_{p_2} + x_{p_2}^2 + y_0^2 - 2y_0 y_{v_2} + y_{v_2}^2 + y_0^2 - 2y_0 y_{p_2} + y_{p_2}^2, \\ \alpha_5 &= 2(x_{v_2}(x_0 - x_{p_2}) + x_{v_0}(-x_0 + x_{p_2}) - (y_0 - y_{v_2})(y_0 - y_{p_2})), \\ \alpha_6 &= 2(x_{p_2}(y_0 - y_{v_2}) + x_{p_0}(-y_0 + y_{v_2}) + (x_{v_0} - x_{v_2})(y_0 - y_{p_2})), \\ \alpha_7 &= x_{v_1}^2 - 2x_{v_1} x_{v_2} + x_{v_2}^2 + x_{p_1}^2 - 2x_{p_1} x_{p_2} + x_{p_2}^2 + y_{v_1}^2 - 2y_{v_1} y_{v_2} + y_{v_2}^2 + y_{p_1}^2 - 2y_{p_1} y_{p_2} + y_{p_2}^2, \\ \alpha_8 &= 2(x_{v_2}(x_{p_1} - x_{p_2}) + x_{v_1}(-x_{p_1} + x_{p_2}) - (y_{v_1} - y_{v_2})(y_{p_1} - y_{p_2})), \\ \alpha_9 &= 2(x_{p_2}(y_{v_1} - y_{v_2}) + x_{p_1}(-y_{v_1} + y_{v_2}) + (x_{v_1} - x_{v_2})(y_{p_1} - y_{p_2})), \\ \alpha_{10} &= 2(x_{v_1} y_{v_0} - x_{v_2} y_{v_0} - x_{v_0} y_{v_1} + x_{v_2} y_{v_1} + x_{v_0} y_{v_2} - x_{v_1} y_{v_2} + x_{p_1} y_{p_0} - x_{p_2} y_{p_0} - x_{p_0} y_{p_1} \\ &\quad + x_{p_2} y_{p_1} + x_{p_0} y_{p_2} - x_{p_1} y_{p_2}), \\ \alpha_{11} &= 2(x_{p_2} y_{v_0} + x_{p_0} y_{v_1} - x_{p_2} y_{v_1} - x_{p_0} y_{v_2} - x_{p_1} y_{v_0} + x_{p_1} y_{v_2} - x_{v_1} y_{p_0} + x_{v_2} y_{p_0} + x_{v_0} y_{p_1} \\ &\quad - x_{v_2} y_{p_1} - x_{v_0} y_{p_2} + x_{v_1} y_{p_2}), \\ \alpha_{12} &= 2(-x_{v_1} x_{p_0} + x_{v_2} x_{p_0} + x_{v_0} x_{p_1} - x_{v_2} x_{p_1} - x_{v_0} x_{p_2} + x_{v_1} x_{p_2} - y_{v_1} y_{p_0} + y_{v_2} y_{p_0} + y_{v_0} y_{p_1} \\ &\quad - y_{v_2} y_{p_1} - y_{v_0} y_{p_2} + y_{v_1} y_{p_2}). \end{aligned}$$

Since  $g(t) = f(\theta) \geq 0$ ,  $\arg \min_t g(t) = \arg \min_t G(t) = g^2(t)$ . To solve  $\min_t G(t)$ , we differentiate  $G(t)$  and take the numerator of  $G'(t)$  as  $P(t)$ . Since  $P(t)$  is a tenth degree polynomial, we use

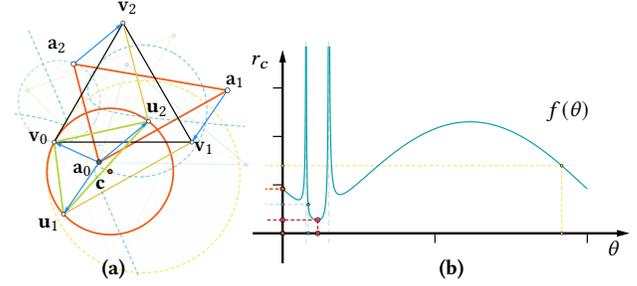


Fig. 3. The acute triangle case. (a)  $\Delta v_0 v_1 v_2$  is a triangle of  $\mathcal{R}$  and  $\Delta a_0 a_1 a_2$  indicates the template triangle  $\Delta p_0 p_1 p_2$  after a rigid transformation  $R$ ,  $\mathbf{b}$ . Quadrilaterals  $a_0 a_1 v_1 u_1$  and  $a_0 a_2 v_2 u_2$  are parallelograms.  $\mathbf{c}$  is the center of the circumcircle of  $\Delta a_0 u_1 u_2$ , which is the same as the MCC of  $\Delta u_0 u_1 u_2$  for the acute triangle. (b) The image of the function  $f(\theta) = r_c$ .

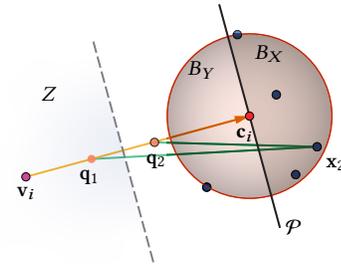


Fig. 4. Blue points in the figure are the templates' vertices corresponding to  $v_i$ , i.e., points in  $S(v_i)$ . The orange circle represents the minimum bounding sphere of  $S(v_i)$ , denoted as  $B_{S(v_i)}$ .  $c_i$  is the center of  $B_{S(v_i)}$ .  $\mathcal{P}$  is the plane perpendicular to the line  $\overline{v_i c_i}$  through the point  $c_i$  dividing the minimum bounding sphere  $B_{S(v_i)}$  into two parts  $B_X$  and  $B_Y$ . The bisecting plane of the edge  $q_1 q_2$  divide the space into two parts, and we denote  $Z$  as the part containing  $q_1$ . When moving  $v_i$  from its start point to  $c_i$ ,  $\max_{s \in S(v_i)} \|s - (v_i)\|$  decreases monotonically and so does the  $\max_{f \in \Omega_i} d_{\text{assembly}}(\mathbf{f})$ .

Jenkins-Traub algorithm [Jenkins and Traub 1970] to find its all ten roots. We first compare the function values at the ten roots to find the minimum and then compute the corresponding  $\theta$  as the result.

### 1.4 Proposition 3

Prop. 3. For each vertex  $v_i$ , the maximum assembly error on its one-ring triangles (denoted as  $\Omega_i$ ) of  $v_i$  is:

$$d(\alpha_i) = \max_{\mathbf{f} \in \Omega_i} d_{\text{assembly}}(\mathbf{f}) = \max_{\mathbf{f} \in \Omega_i} \min_{t \in \mathcal{T}} d_{\max}(\mathbf{f}, t^{\delta_j}) \quad (1)$$

where  $0 \leq \alpha_i \leq 1$  is the step size. Then,  $d(\alpha_i)$  monotonically decreases with respect to  $\alpha_i$ .

PROOF. Let  $\mathcal{P}$  be the plane perpendicular to the line  $\overline{v_i c_i}$  through the point  $c_i$  (Fig. 4).  $S(v_i)$  consists of the templates' vertices corresponding to  $v_i$ .  $\mathcal{P}$  divides the minimum bounding sphere  $B_{S(v_i)}$  into two parts  $B_X$  and  $B_Y$ . Let  $B_X$  be the part far away from  $v_i$  and  $B'_X = B_X \cup \mathcal{P}$ . Then, the point set  $X = S(v_i) \cap B'_X$  are not empty; otherwise, there is a bounding sphere having a smaller radius than  $B_{S(v_i)}$ , which contradicts the assertion that  $B_{S(v_i)}$  is the minimum bounding sphere. Let  $v'_i = v_i + \alpha_i \mathbf{d}$ , we define:

$$\delta(\alpha_i) = d_{S(v_i)}(v'_i) = \max_{s \in S(v_i)} \|s - (v_i + \alpha_i \mathbf{d})\|, \quad (2)$$

where  $d_S(\mathbf{v}) = \max_{s \in S} \|\mathbf{v} - \mathbf{s}\|_2$  is defined as the distance from a point  $\mathbf{v}$  to a point set  $S$ .

$\delta(\alpha_i)$  is monotonically decreasing, and we prove it by contradiction. According to the fact that  $\delta(\alpha_i) = d_S(\mathbf{v}'_i) = d_X(\mathbf{v}'_i) \geq d_Y(\mathbf{v}'_i)$ , we only need to focus on the set  $X$  and  $d_X(\mathbf{v}'_i)$ . Suppose that  $\delta(\alpha_i), \alpha_i \in [0, 1]$  does not monotonically decrease, then there are two different values  $\alpha_i^1 < \alpha_i^2$ , s.t.  $\delta(\alpha_i^1) < \delta(\alpha_i^2)$ . Namely,  $\exists \mathbf{q}_1, \mathbf{q}_2 \in \overline{\mathbf{v}_i \mathbf{c}_i}, \mathbf{q}_1 = \mathbf{v}_i + \alpha_i^1 \mathbf{d}, \mathbf{q}_2 = \mathbf{v}_i + \alpha_i^2 \mathbf{d}$  and  $\alpha_i^1 < \alpha_i^2$ , s.t.  $d_X(\mathbf{q}_1) < d_X(\mathbf{q}_2)$ . Let  $\mathbf{x}_2 \in X$  be the point where  $d_X(\mathbf{q}_2)$  is obtained, i.e.,  $d_X(\mathbf{q}_2) = \|\mathbf{q}_2 - \mathbf{x}_2\|$ . Then,  $\|\mathbf{q}_1 - \mathbf{x}_2\| \leq d_X(\mathbf{q}_1) < d_X(\mathbf{q}_2) = \|\mathbf{q}_2 - \mathbf{x}_2\|$ . The bisecting plane of the edge  $\mathbf{q}_1 \mathbf{q}_2$  divide the space into two parts, and we denote  $Z$  as the part containing  $\mathbf{q}_1$ . We have  $Z \cap X = \emptyset$ , and since  $\mathbf{x}_2 \in X, \mathbf{x}_2 \notin Z$ . Thus,  $\|\mathbf{q}_1 - \mathbf{x}_2\| \geq \|\mathbf{q}_2 - \mathbf{x}_2\|$ , and the contradiction arises.

Then,  $\forall \alpha_i \in [0, 1]$ , we have:

$$d(\alpha_i) = \max_{\mathbf{f} \in \Omega_i} d_{\text{assembly}}(\mathbf{f}) = \max_{\mathbf{f} \in \Omega_i} \min_{\substack{\mathbf{t} \in \mathcal{T} \\ j \in \{1, \dots, 6\}}} d_{\max}(\mathbf{f}, \mathbf{t}^{\phi_j}) \geq \delta(\alpha_i). \quad (3)$$

$d(\alpha_i) = \max_{\mathbf{f} \in \Omega_i} d_{\text{assembly}}(\mathbf{f})$  must be obtained at one vertex on certain triangle  $\mathbf{f} \in \Omega_i$ , denoted as  $\mathbf{v}^*$ . If  $\mathbf{v}^* = \mathbf{v}'_i$ , then  $d(\alpha_i) =$

$\delta(\alpha_i)$ ; otherwise,  $d(\alpha_i)$  is obtained at another point, indicating that  $d(\alpha_i)$  does not change as  $\alpha_i$  updates. More specifically, if  $\exists \alpha'_i$ , s.t.  $d(\alpha'_i) > \delta(\alpha'_i)$ , then  $d(\alpha_i)$  is constant  $\forall \alpha_i \in (\alpha'_i, 1)$ ; otherwise  $d(\alpha_i) = \delta(\alpha_i)$ . Hence, since  $\delta(\alpha_i)$  is monotonically decreasing,  $d(\alpha_i)$  is also monotonically decreasing.  $\square$

### 1.5 Proposition 4

*Prop. 4.* Given two triangles  $A = \Delta \mathbf{a}_0 \mathbf{a}_1 \mathbf{a}_2$  and  $B = \Delta \mathbf{b}_0 \mathbf{b}_1 \mathbf{b}_2$ , let  $d_H(A, B)$  be the two-sided Hausdorff distance between two triangles and  $\xi = (\|\mathbf{b}_0 - \mathbf{a}_0\|_2, \|\mathbf{b}_1 - \mathbf{a}_1\|_2, \|\mathbf{b}_2 - \mathbf{a}_2\|_2)$ . Then,

$$d_H(A, B) \leq \|\xi\|_\infty \leq \|\xi\|_2.$$

**PROOF.** Since  $\forall i \in \{0, 1, 2\}, d_H(A, B) \leq \|\mathbf{b}_i - \mathbf{a}_i\|_2 \leq \|\xi\|_2, d_H(A, B) \leq \|\xi\|_\infty \leq \|\xi\|_2. \square$

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